# MetaGrad: Faster Convergence **Without Curvature** in Online Convex Optimization



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- Online Convex Optimization
- Learning the Learning rate
- Second-order (variance) bounds (individual sequence)
- Fast rates without curvature



#### **Online Convex Optimization**

A New Type of Guarantee

Fast Rates

MetaGrad Algorithm

#### Fundamental Learning Model: Online Convex Optimization

- ln round  $t = 1, 2, \ldots$ 
  - Learner predicts  $w_t$  (from unit ball)
  - Encounter convex loss function  $f_t(u): \mathbb{R}^d o \mathbb{R}$



- Learner
  - observes gradient  $g_t \coloneqq 
    abla f_t(w_t)$  (from unit ball)
  - incurs loss  $f_t(w_t)$

#### Fundamental Learning Model: Online Convex Optimization

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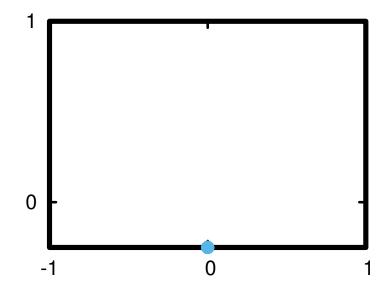


- Learner
  - observes gradient  ${m g}_t\coloneqq 
    abla f_t({m w}_t)$  (from unit ball)
  - incurs loss f<sub>t</sub>(w<sub>t</sub>)
- The goal is to have small regret

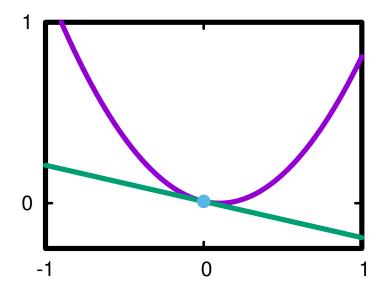
$$R_T^u := \underbrace{\sum_{t=1}^T f_t(w_t)}_{\text{Learner}} - \underbrace{\sum_{t=1}^T f_t(u)}_{\text{Point } u}$$

with respect to every point u.

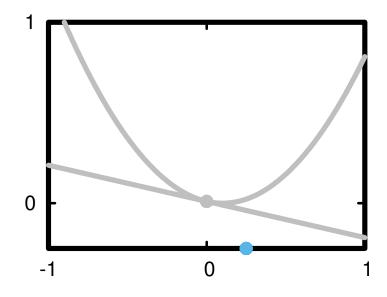
Round 1: Learner plays  $w_1 = 0$ 



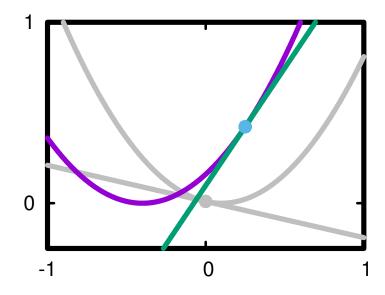
Round 1: Learner incurs  $f_1(w_1)$  and sees  $g_1 = \nabla f_1(w_1)$ 



Round 2: Learner plays  $w_1 = 1/4$ 

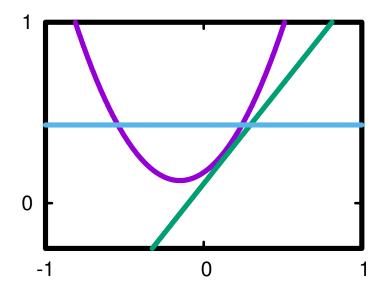


Round 2: Learner incurs  $f_2(w_2)$  and sees  $g_2 = \nabla f_2(w_2)$ 



. . .

Evaluate Learner using **regret**:  $R_T^u = \sum_{t=1}^T (f_t(w_t) - f_t(u))$ 



State of the Art



Online gradient descent

$$oldsymbol{w}_{t+1} \;=\; oldsymbol{w}_t - oldsymbol{\eta} oldsymbol{g}_t$$

recall  $\boldsymbol{g}_t = 
abla f_t(\boldsymbol{w}_t)$ 

State of the Art



Online gradient descent

$$oldsymbol{w}_{t+1} \;=\; oldsymbol{w}_t - oldsymbol{\eta} oldsymbol{g}_t$$

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abla f_t(\boldsymbol{w}_t)$ 

OGD bound: After T rounds,

$$R_T^{\boldsymbol{u}} \leq O\left(\sqrt{\sum_{t=1}^T \lVert \boldsymbol{g}_t 
Vert^2}
ight)$$

for all  $\boldsymbol{u}$  with  $\|\boldsymbol{u}\| \leq 1$ .

## Bounds Reveal Our Dearest Hopes

Always have worst-case guarantee

$$R_T^{\boldsymbol{u}} \leq O\left(\sqrt{\sum_{t=1}^T \|\boldsymbol{g}_t\|^2}
ight) \leq O(\sqrt{T}).$$

#### Yet bound says we might get lucky

For smooth functions  $f_t$  with common optimum  $m{u}^*$ , as  $m{w}_t o m{u}^*$ , we have  $m{g}_t o m{0}$ , and

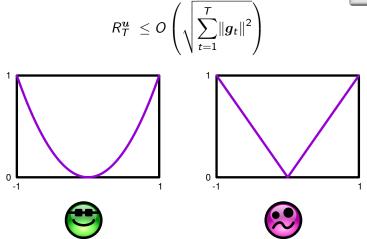
$$\sqrt{\sum_{t=1}^{\mathcal{T}} \lVert oldsymbol{g}_t 
Vert^2} \ \ll \ \sqrt{\mathcal{T}}$$

grows much slower than  $\sqrt{T}$ .



## What We Hope Happens





## Can We Do Better?

No in general: matching lower bound.

 $R_T^u \geq \Omega(\sqrt{T})$ 



Yes, with curvature:

 $R_T^u \leq O(\ln T)$ 

• Strongly convex: 
$$I \preceq 
abla^2 f(u)$$
, e.g.

$$f_t(\boldsymbol{u}) = \|\boldsymbol{u} - \boldsymbol{y}_t\|^2$$

 $\Rightarrow$  gradient descent with small  $\eta$ 

• Exp-concave:  $\nabla f(u) \nabla f(u)^{\intercal} \preceq \nabla^2 f(u)$ , e.g.

$$f_t(u) ~=~ -\ln\left(1+m{y}_t^{\intercal}u
ight)$$

 $\Rightarrow$  Online Newton Step



But do we really need curvature?

This talk: no, **stability** is enough.

New algorithm MetaGrad: Separate learning rate  $\eta$  for each point u



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#### Refined Bound

Recall bound for gradient descent:

$$R_T^{\boldsymbol{u}} \leq O\left(\sqrt{\sum_{t=1}^T \lVert \boldsymbol{g}_t 
Vert^2}
ight)$$

New bound for MetaGrad:

$$R_T^{\boldsymbol{u}} \leq O\left(\sqrt{V_T^{\boldsymbol{u}} d \ln T}\right) \quad \text{where} \quad V_T^{\boldsymbol{u}} \coloneqq \sum_{t=1}^T \left((\boldsymbol{w}_t - \boldsymbol{u})^{\mathsf{T}} \boldsymbol{g}_t\right)^2$$

Data-dependent. Whoa! Ouroboric.

Always improvement:

$$\left((\boldsymbol{w}_t - \boldsymbol{u})^{\intercal} \boldsymbol{g}_t
ight)^2 \leq \|\boldsymbol{w}_t - \boldsymbol{u}\|^2 \|\boldsymbol{g}_t\|^2$$



Now What We Hope Happens

$$R_T^{\boldsymbol{u}} \leq O\left(\sqrt{\sum_{t=1}^T ((\boldsymbol{w}_t - \boldsymbol{u})^{\mathsf{T}}\boldsymbol{g}_t)^2}\right)$$



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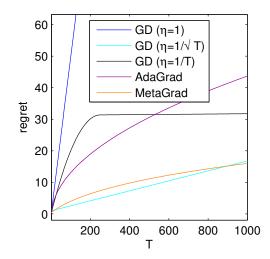
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#### Does it Really Work?

Offline optimization (fixed function):

$$f_t(u) = |u-1/4|$$





The "Fast Rates" Pipeline



#### Combine

refined individual-sequence regret bound

$$R_T^u \leq \sqrt{V_T^u d \ln T} \quad \forall u$$

Special-purpose argument that for best  $u^*$ 

$$\boxed{V_T^{\boldsymbol{u}^*} \leq R_T^{\boldsymbol{u}^*}}$$

Profit!

$$R_T^{\boldsymbol{u}^*} \leq \sqrt{R_T^{\boldsymbol{u}^*} d \ln T}$$
 so  $R_T^{\boldsymbol{u}^*} \leq d \ln T$ 

# Significant Improvement: Fixed Function

Any fixed  $f_t(u) = f(u)$ . Let  $u^* = \arg \min_u f(u)$  be the offline minimiser.

Crux: 
$$(\boldsymbol{w}_t - \boldsymbol{u}^*)^{\intercal} \boldsymbol{g}_t \in [0, 2].$$

Now from the regret bound

$$R_T^{oldsymbol{u}^*} \leq \sum_{t=1}^T (oldsymbol{w}_t - oldsymbol{u}^*)^\intercal oldsymbol{g}_t \leq \sqrt{V_T^{oldsymbol{u}^*} d \ln T}$$

and special-purpose observation

$$V_T^{{m u}^*} = \sum_{t=1}^T \left( ({m u}^* - {m w}_t)^{\mathsf{T}} {m g}_t 
ight)^2 \le 2 \sum_{t=1}^T ({m w}_t - {m u}^*)^{\mathsf{T}} {m g}_t$$

we can solve for  $V_T^{{m u}^*}$  to find  $V_T^{{m u}^*} \leq d \ln T$  and hence

$$\mathsf{R}_T^{oldsymbol{u}^*}~\leq~\sqrt{2}d\,{
m ln}~7$$

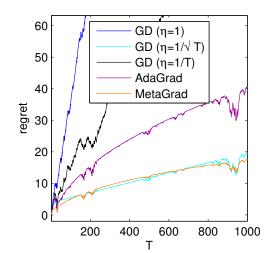


## Does It Really Actually Work?

Stochastic optimization:

$$f_t(u) = |u - x_t|$$

where  $x_t = \pm \frac{1}{2}$  i.i.d. with probability 0.4 and 0.6.



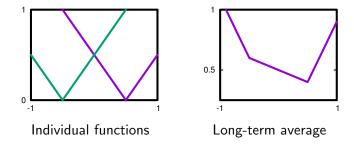


## What's Going On, Really?

Stochastic optimization:

$$f_t(u) = |u - x_t|$$

where  $x_t = \pm \frac{1}{2}$  i.i.d. with probability 0.4 and 0.6.



Stable minimum easy to converge to

Significant Improvement: Stochastic Case Consider i.i.d.

$$f \sim \mathbb{P}$$
 with  $u^* = rgmin \mathop{\mathbb{E}}_{u}[f(u)]$ 

Condition: there is a c > 0 such that

 $\nabla \forall m{w}: (m{w}-m{u}^*)^\intercal \, \mathbb{E}\left[ 
abla f(m{w}) 
abla f(m{w})^\intercal 
ight] (m{w}-m{u}^*) \ \leq \ m{c}(m{w}-m{u}^*)^\intercal \, \mathbb{E}\left[ 
abla f(m{w}) 
ight]$ 

Now from the special-case condition

$$\mathbb{E}[V_T^{oldsymbol{u}^*}] \ \le \ c \, \mathbb{E}\left[\sum_{t=1}^T (oldsymbol{w}_t - oldsymbol{u}^*)^\intercal 
abla f(oldsymbol{w}_t)
ight]$$

and by the generic regret bound, in expectation,

$$\mathbb{E}[R_T^{\boldsymbol{u}^*}] \leq \mathbb{E}\left[\sum_{t=1}^T (\boldsymbol{w}_t - \boldsymbol{u}^*)^{\mathsf{T}} \nabla f(\boldsymbol{w}_t)\right] \leq \mathbb{E}\left[\sqrt{V_T^{\boldsymbol{u}^*} d \ln T}\right]$$
  
and by Jensen's inequality  $\mathbb{E}\left[\sqrt{V_T^{\boldsymbol{u}^*}}\right] \leq \sqrt{\mathbb{E}\left[V_T^{\boldsymbol{u}^*}\right]}$ , so that  
 $\mathbb{E}[R_T^{\boldsymbol{u}^*}] \leq \sqrt{c} d \ln T$ 



#### Outline

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# Our Approach in a Nutshell



- 1. Replace actual loss  $f_t(u)$  by surrogate loss  $\ell_t^{\eta}(u)$ 
  - parametrised by learning rate  $\eta$
  - $\blacktriangleright$  exp-concave in u
  - So can get good bound for surrogate regret
- 2. Exponentially spaced grid  $\eta_1, \eta_2, \ldots, \eta_{\log(T)}$   $(\eta_i = 2^{-i})$ .
- 3. Off-the-shelf exp-concave **Slave** for grid point  $\eta_i$  predicts

$$w_1^{\eta_i}, w_2^{\eta_i}, \ldots$$

4. At each round t, Master aggregates  $w_t^{\eta_1}, w_t^{\eta_2}, \ldots$  into  $w_t$ .

## Surrogate Loss



Real loss

$$f_t(\boldsymbol{u}) ~\leq~ f_t(\boldsymbol{w}_t) + (\boldsymbol{u} - \boldsymbol{w}_t)^{\intercal} \boldsymbol{g}_t$$

Surrogate loss

$$\ell^\eta_t(u) \ \coloneqq \ \eta(u-w_t)^{\intercal} g_t + (\eta(u-w_t)^{\intercal} g_t)^2$$

Exp-concave! In particular:

$$e^{-\ell_t^{\boldsymbol{\eta}}(\boldsymbol{u})} \leq 1 + \boldsymbol{\eta}(\boldsymbol{w}_t - \boldsymbol{u})^{\mathsf{T}} \boldsymbol{g}_t.$$

Excellent bound  $O(\ln T)$  for wrong loss.

#### MetaGrad Slave

 $\eta$ -Slave (variant of Online Newton Step) predicts

$$w_{t+1}^{\eta} ~=~ w_t^{\eta} - \eta \Sigma_{t+1}^{\eta} g_t$$

where the covariance matrix is given by

$$\boldsymbol{\Sigma}_{t+1}^{\boldsymbol{\eta}} = \left(\frac{1}{4}\boldsymbol{I} + 2\boldsymbol{\eta}^2\sum_{s=1}^{t}\boldsymbol{g}_s\boldsymbol{g}_s^{\mathsf{T}}\right)^{-1}$$

 $\eta$ -Slave guarantees

$$\begin{split} \sum_{t=1}^T \bigl(\ell_t^{\boldsymbol{\eta}}(\boldsymbol{w}_t^{\boldsymbol{\eta}}) - \ell_t^{\boldsymbol{\eta}}(\boldsymbol{u})\bigr) &\leq \frac{1}{8} \|\boldsymbol{u}\|^2 + \frac{1}{2} \ln \det \left(\boldsymbol{I} + 8\boldsymbol{\eta}^2 \sum_{t=1}^T \boldsymbol{g}_t \boldsymbol{g}_t^{\mathsf{T}}\right) \\ &\leq O(d \ln \mathcal{T}) \quad \forall \boldsymbol{u} \end{split}$$



#### MetaGrad Master



Input: Grid points  $\eta_i = 2^{-i}$  with weights  $\pi_i = \frac{1}{i(i+1)}$ . Goal: aggregate  $w_t^{\eta_1}, w_t^{\eta_2}, \ldots$ 

Idea: Potential

$$\Phi_t := \sum_i \pi_i e^{-\sum_{s=1}^t \ell_s^{\eta_i}(w_s^{\eta_i})}.$$

Two steps:

- Find predictions  $w_t$  that ensure  $1 \ge \Phi_1 \ge \Phi_2 \ge \ldots$
- Derive regret bound from  $1 \ge \Phi_T$ .

#### MetaGrad Master, Potential Decreases



#### Tilted exponentially weighted average

$$w_{t+1} = \frac{\sum_{i} \pi_{i} e^{-\sum_{s=1}^{t} \ell_{s}^{\eta_{i}}(w_{s}^{\eta_{i}})} \eta_{i} w_{t+1}^{\eta_{i}}}{\sum_{i} \pi_{i} e^{-\sum_{s=1}^{t} \ell_{s}^{\eta_{i}}(w_{s}^{\eta_{i}})} \eta_{i}}$$

ensures potential shrinks:

$$\begin{split} \Phi_{t+1} - \Phi_t &= \sum_{i} \pi_i e^{-\sum_{s=1}^t \ell_s^{\eta_i}(w_s^{\eta_i})} \left( e^{-\ell_{t+1}^{\eta_i}(w_{t+1}^{\eta_i})} - 1 \right) \\ & \leq \sum_{i} \pi_i e^{-\sum_{s=1}^t \ell_s^{\eta_i}(w_s^{\eta_i})} \eta_i (w_{t+1} - w_{t+1}^{\eta_i})^{\mathsf{T}} g_{t+1} \overset{\text{weights}}{=} 0 \end{split}$$

and hence  $\Phi_t \leq 1$ .

#### MetaGrad Master, Small Potential is Good



The Master achieves for all *t*:

$$1 \geq \Phi_t = \sum_{i} \pi_i e^{-\sum_{s=1}^t \ell_s^{\eta_i}(w_s^{\eta_i})}.$$

It follows that

$$\sum_{t=1}^T \left( 0 - \ell_t^{oldsymbol{\eta}_i}(oldsymbol{w}_t^{oldsymbol{\eta}_i}) 
ight) \ \le \ -\ln \pi_i \qquad orall i \ ext{in grid}$$

(Master has zero surrogate loss)

## MetaGrad Analysis

Now combine the Master and Slave guarantee. For **each** grid point  $\eta$  and comparator u

$$\sum_{t=1}^{T} \left( 0 - \ell_t^{\eta}(\boldsymbol{w}_t^{\eta}) \right) \leq -\ln \pi_i \leq \ln \ln 7$$
$$\sum_{t=1}^{T} \left( \ell_t^{\eta}(\boldsymbol{w}_t^{\eta}) - \ell_t^{\eta}(\boldsymbol{u}) \right) \leq O(d \ln 7)$$

SO

$$\sum_{t=1}^T (0-\ell_t^{\eta}(u)) \leq O(d \ln T).$$

Unpacking  $\ell^\eta_t(u) = \eta(u-w_t)^\intercal g_t + (\eta(u-w_t)^\intercal g_t)^2$  yields

 $\eta R_T^u \leq \eta^2 V_T^u + O(d \ln T).$ 



MetaGrad Analysis (ctd.)

Reorganise the bound to:

$$R_T^{oldsymbol{u}} \ \le \ oldsymbol{\eta} V_T^{oldsymbol{u}} + rac{O(d \ln T)}{\eta}$$

Now pick the best grid point

$$\hat{\eta} = \sqrt{\frac{O(d \ln T)}{V_T^u}}$$

to find

$$R_T^u \leq O\left(\sqrt{V_T^u d \ln T}\right)$$

of course we need a grid point close to  $\hat{\eta}$  and we need to deal with off-grid  $\hat{\eta} \gg 1$  and  $\hat{\eta} \ll \frac{1}{\sqrt{T}}$ .



# MetaGrad Outlook

- Run-time O(d<sup>2</sup>) per round
- Projections (avoid O(d<sup>3</sup>) per round!)
- We design and analyze two versions of Slave
  - Full covariance (quadratic)
  - Diagonal approximation (linear)
- Very welcome to discuss further



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Learn more:

- Paper submitted to COLT 2016, preprint available
- Code is available



http://bitbucket.org/wmkoolen/metagrad

Experiments coming soon.

http://blog.wouterkoolen.info





# Low regret through stability, even without curvature.

- New MetaGrad algorithm.
- Hierarchical Master-Slave construction.
- Learns the learning rate.
- Refined (adaptive) regret bound.
- Stochastic condition for logarithmic regret (fast rates)

# Thank you!