## A Closer Look at Adaptive Regret

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Predictor



Expert Expert Expert

Nature



TFI

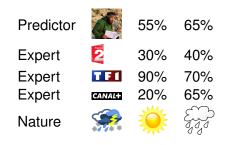
PredictorImage: Second sec

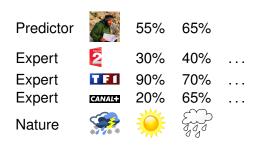


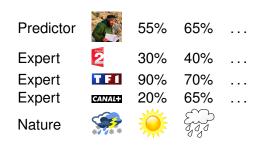


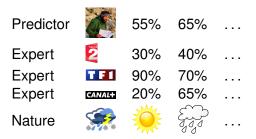






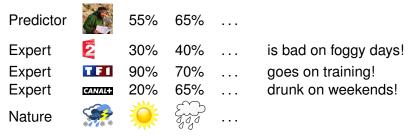


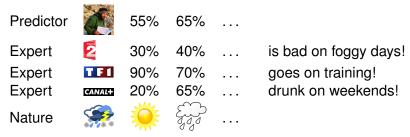






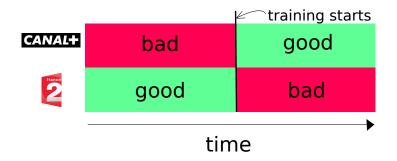




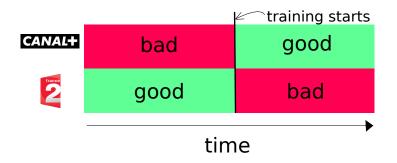


Goal: close to the best expert overall (solution: AA) Adaptive goal: close to the best expert on every time interval

# Example continued



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#### Non-adaptive predictor would lose trust in the first guy.

 Blowing up the set of experts to compete with virtual sleeping experts [Adamskiy, Koolen, Chernov and Vovk 2012]

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#### Our results

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## Our results

- Figured out the worst-case adaptive regret of Fixed Share
- Proved the optimality of Fixed Share every algorithm is dominated by a Fixed Share (i.t.o. worst-case adaptive regret)







for t = 1, 2, ... do

Learner announces probability vector  $\boldsymbol{w}_t \in \Delta_N$ Reality announces loss vector  $\boldsymbol{\ell}_t \in [-\infty, \infty]^N$ Learner suffers loss  $\boldsymbol{\ell}_t := -\ln \sum_n \boldsymbol{w}_t^n \boldsymbol{e}^{-\ell_t^n}$ end for

- Goal: On every time interval [*t*<sub>1</sub>, *t*<sub>2</sub>] we want to be not much worse than the best expert on that interval.
- We are interested in small adaptive regret

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## Definition

The adaptive regret of the algorithm on the interval  $[t_1, t_2]$  is the loss of the algorithm there minus the lost of the best expert there:

$$R_{[t_1,t_2]} := L_{[t_1,t_2]} - \min_j L^j_{[t_1,t_2]}$$

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 We get worst-case adaptive regret by maximising over the choice of expert losses. Aggregating Algorithm [Vovk 1990] updates weights as:

$$w_{t+1}^n := \frac{w_t^n e^{-\ell_t^n}}{\sum_n w_t^n e^{-\ell_t^n}}.$$

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Fixed Share (family) is defined by a sequence of "switching rates"  $\alpha_t$ . The weight update is

$$w_{t+1}^n := \frac{\alpha_{t+1}}{N-1} + \left(1 - \frac{N}{N-1}\alpha_{t+1}\right) \frac{w_t^n e^{-\ell_t^n}}{\sum_n w_t^n e^{-\ell_t^n}}.$$

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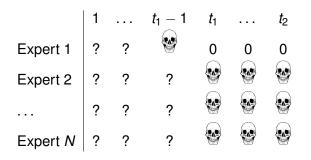
Adaptivity hides in the first term.







We proved that the worst case data for Fixed Share looks like this:



where denotes infinite loss, 0 – zero loss and '?' – losses that don't matter.

Knowing the worst-case data, we can plug it in and calculate the regret:

#### Theorem

The worst-case adaptive regret of Fixed Share with N experts on interval  $[t_1, t_2]$  equals

$$-\ln\left(\frac{\alpha_{t_1}}{N-1}\prod_{t=t_1+1}^{t_2}(1-\alpha_t)\right)$$

• Classic Fixed Share ( $\alpha_t = const$ ):

$$\ln(N-1) - \ln \alpha - (t_2 - t_1) \ln(1 - \alpha)$$
 for  $t_1 > 1$ , and   
  $\ln N - (t_2 - 1) \ln(1 - \alpha)$  for  $t_1 = 1$ .

Linear in length  $t_2 - t_1$ 

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# Linear in length $t_2 - t_1$

• Slowly decreasing  $\alpha_t = 1/t$  leads to regret of

$$\begin{aligned} &\ln(N-1) + \ln t_2 & \text{for } t_1 > 1, \text{ and} \\ &\ln N + \ln t_2 & \text{for } t_1 = 1. \end{aligned}$$

#### Logarithmic in end-time t<sub>2</sub>

# • Quickly decreasing switching rate. If we set $\alpha_t = t^{-2}$ we have the upper bound for regret

```
\ln N + 2 \ln t_1 + \ln 2.
```

#### Logarithmic in start-time $t_1$

For  $t_1 = 1$  this is very close to classical AA regret!

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Time	Interval 1				Interval 2				Interval 3				
Expert 1	0	0		0									
Expert 2	•				0	0		0					
Expert 3	**								0	0		0	

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And the tracking bound can be recovered!

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#### Theorem

The worst-case adaptive regret of any algorithm is dominated by that of a Fixed Share.

# Theorem (Lower bound)

Let N be the number of experts. If  $\phi(t_1, t_2)$  is the worst-case adaptive regret of an algorithm, then

$$\begin{split} \phi(t,t) &\geq \ln N & \text{for all } 1 \leq t \text{ and} \\ \phi(t_1,t_2) &\geq \phi(t_1,t_1) - \sum_{t=t_1+1}^{t_2} \ln \left(1 - (N-1)e^{-\phi(t,t)}\right) & \text{for all } 1 \leq t_1 < t_2. \end{split}$$

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Fixed Share with

$$\alpha_t = (N-1)e^{-\phi(t,t)}$$

will have the second line with equality.

- We studied two intuitive methods to obtain adaptive algorithms.
- They turned out to be Fixed Share.
- We studied the worst-case Adaptive Regret of Fixed Share.
- We showed that Fixed Share is optimal.

Thank you!