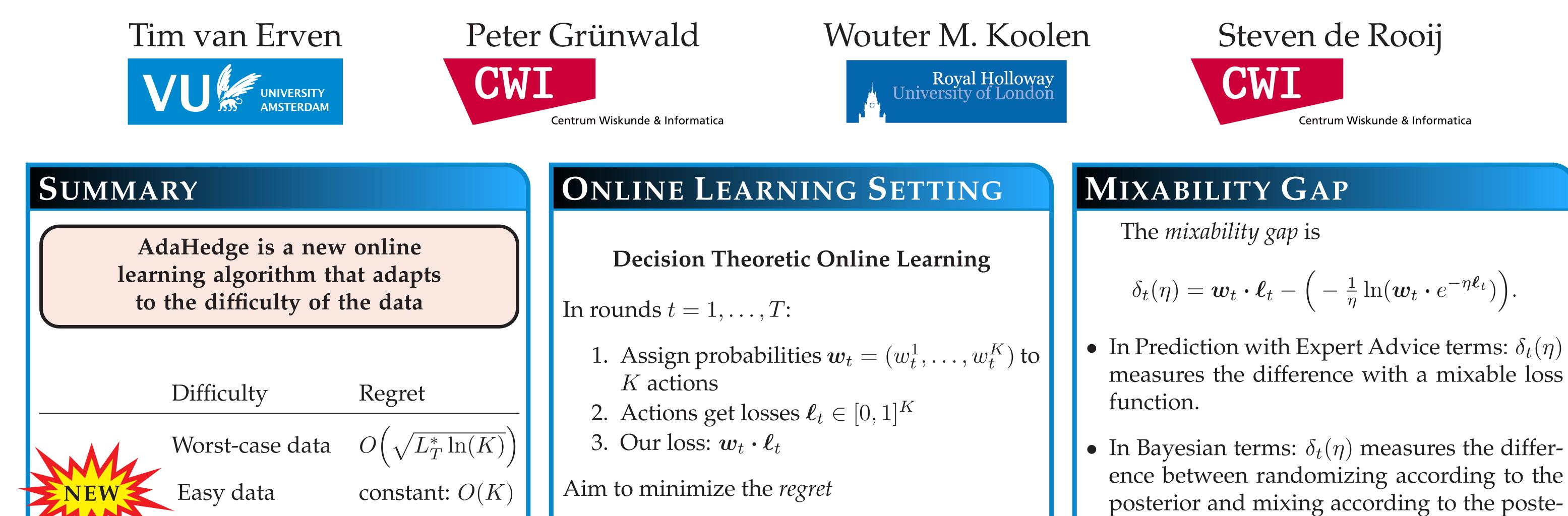
# **Adaptive Hedge**





**Key Ideas** 

- Bounds on the *mixability gap* (see top-right panel) play a crucial role in previous analyses of the Hedge algorithm.
- We only bound the mixability gap in the analysis, but not in the algorithm!
- On easy data, the probabilities output by Hedge converge on a single action. In this case we improve the standard bounds.
- Example: if one action is always better than all others.

$$R(T) = \sum_{t=1}^{T} \boldsymbol{w}_t \cdot \boldsymbol{\ell}_t - L_T^*,$$

where  $L_T^* = \min_k \sum_{t=1}^T \ell_t^k$  is the loss of the best action in hindsight.

# HEDGE

• Hedge predicts with exponential weights:

$$w_t^k \propto \exp\left(-\eta \sum_{s=1}^{t-1} \ell_s^k\right).$$

• Its performance depends strongly on the *learning rate*  $\eta > 0$ .

rior.

# ADAHEDGE

- Tune  $\eta$  optimally for a budget  $b(\eta)$  on the cumulative mixability gap  $\Delta_T(\eta) = \sum_{t=1}^T \delta_t(\eta)$
- Increase the budget using the doubling trick.

### Algorithm

- 1. Start with  $\eta = 1$
- 2. Run a new instance of Hedge with learning rate  $\eta$  until  $\Delta_T(\eta)$  exceeds budget

$$b(\eta) = \left(\frac{1}{\eta} + \frac{1}{e-1}\right) \ln(K)$$
  
3. Set  $\eta \leftarrow \eta/2$  and goto 2.

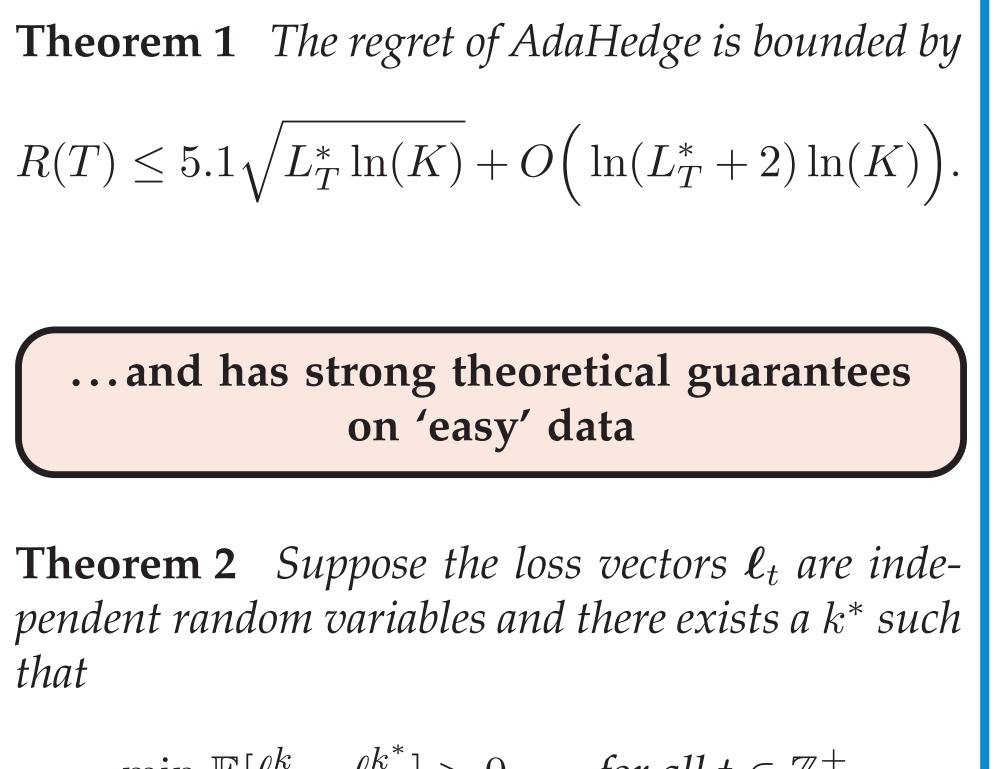
## **THEORETICAL RESULTS**

AdaHedge is worst-case optimal...

## EXPERIMENTS

Simulation Study on 'Easy' Data

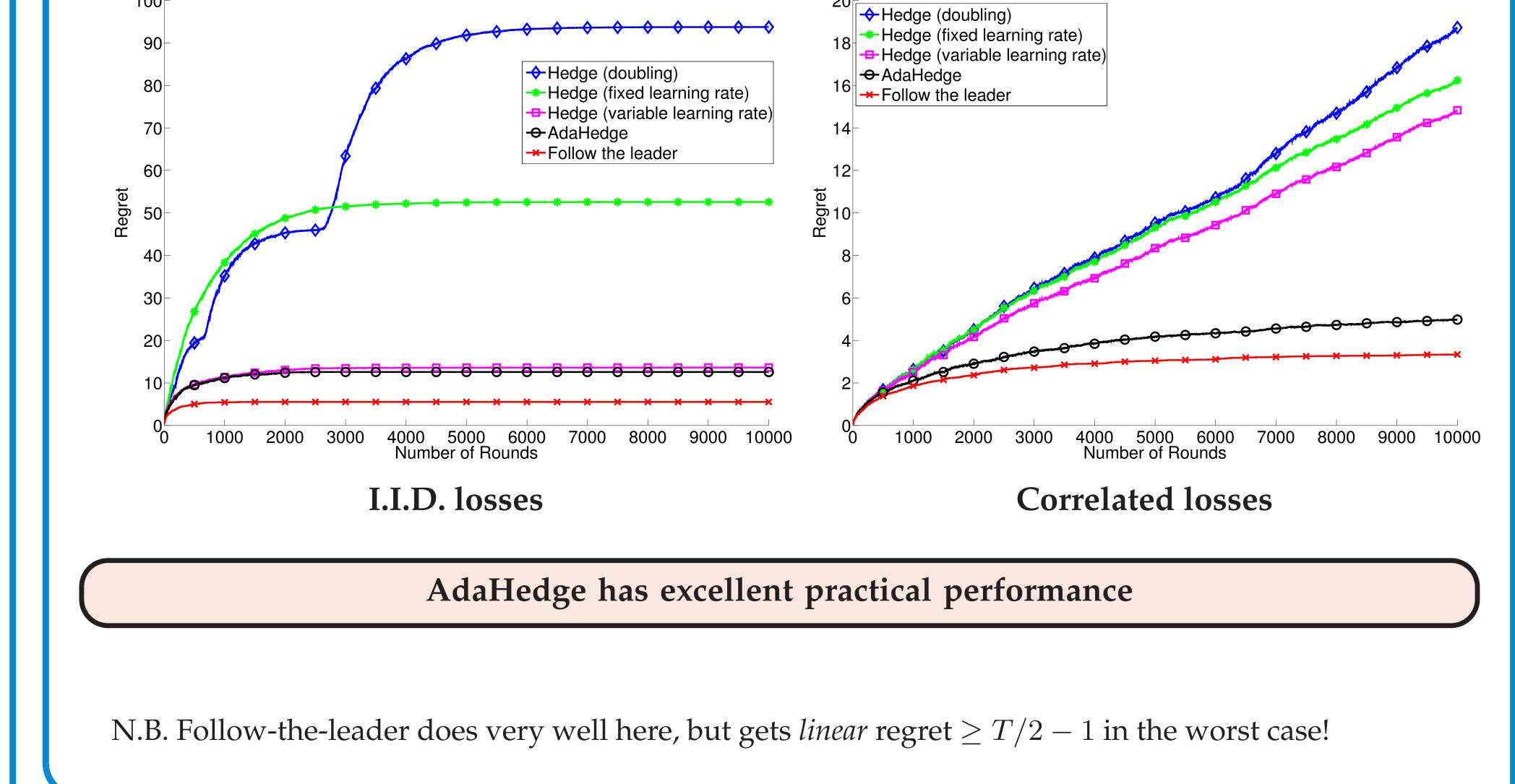
100



 $\min_{k \neq k^*} \mathbb{E}[\ell_t^k - \ell_t^{k^*}] > 0 \quad \text{for all } t \in \mathbb{Z}^+.$ 

Then with probability at least  $1-\delta$  the regret of Ada-*Hedge is bounded by a constant:* 

 $R(T) = O\left(K + \log(1/\delta)\right).$ 



#### **CURRENT WORK**

# **PROOF TECHNIQUES**

Everyone bounds the mixability gap  $\delta_t$ .

**Standard Analysis** 

• Optimize  $\eta$  after bounding  $\delta_t(\eta) \leq \eta/8$ .

**Our Approach** 

- Optimize  $\eta$  before bounding!
- If the posterior probabilities  $w_t$  converge on a single action, the mixability gap goes to 0!

 $\delta_t(\eta) \le (e-2)\eta \left(1 - \max_{k} w_t^k\right) \qquad (0 < \eta \le 1)$ 

Avoid the Doubling Trick

- Better performance in practice
- Still very clean analysis
- Improved the worst-case bound to

 $R(T) \le 2\sqrt{\frac{L_T^*(T - L_T^*)}{T}}\ln(K) + \frac{8}{3}\ln(K) + 2.$ 

#### Weaker Conditions for Easy Data

• Guarantee regret bounded by the best *regret* of AdaHedge and Follow-the-Leader, up to a small constant factor.

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