

Comprendre le monde, construire l'avenir®

Follow the leader if you can, Hedge if you must

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NIPS, 2013

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Outline

- Follow-the-Leader:
 - works well for `easy' data: few leader changes, i.i.d.
 - but not robust to worst-case data
- Exponential weights with simple tuning:
 - robust, but does not exploit easy data
- Second-order bounds:
 - robust against worst case + can exploit i.i.d. data
 - but do not exploit few leader changes in general
- FlipFlop: robust + as good as FTL

Sequential Prediction with Expert Advice

- *K* experts sequentially predict data x_1, x_2, \ldots
- Goal: predict (almost) as well as the best expert on average
- Applications:
 - online convex optimization
 - predicting electricity consumption
 - predicting air pollution levels
 - spam detection

Set-up: Repeated Game

• Every round $t = 1, \ldots, T$:

- 1. Predict probability distribution $w_t = (w_{t,1}, \dots, w_{t,K})$ on experts
- 2. Observe expert losses $\ell_t = (\ell_{t,1}, \dots, \ell_{t,K}) \in [0, 1]^K$ 3. Our loss is $w_t \cdot \ell_t = \sum_k w_{t,k} \ell_{t,k}$

Goal: minimize *regret*

 \mathbf{T}

$$\sum_{t=1}^{I} w_t \cdot \ell_t - L^* \quad \text{where} \quad L^* = \min_k \sum_{t=1}^{I} \ell_{t,k}$$

Loss of the best expert

Follow-the-Leader

 Deterministically choose the expert that has predicted best in the past:

$$w_{t,k^*} = 1$$
 where $k^* = \arg\min_k \sum_{s=1}^{t-1} \ell_{t,k}$

• Equivalently:

$$w_t = \underset{w}{\operatorname{arg\,min}} \mathbb{E}_{k \sim w} \left[\sum_{s=1}^{t-1} \ell_{t,k} \right]$$

FTL: the Good News

- Regret bounded by nr of leader changes
- Proof sketch:
 - If the leader does not change, our loss is the same as the loss of the leader, so the regret stays the same
 - If the leader does change, our regret increases at most by 1 (range of losses)

 Works well for i.i.d. losses, because the leader changes only finitely many times w.h.p.





 4 experts with Bernoulli 0.1, 0.2, 0.3, 0.4 losses

FTL Worst-case Losses



Exponential Weights



$$w_t = \underset{w}{\operatorname{arg\,min}} \mathbb{E}_{k \sim w} \left[\sum_{s=1}^{t-1} \ell_{t,k} \right] + \frac{1}{\eta} D(w \| u)$$

• $\eta \rightarrow \infty$: recover FTL (aggressive learning)

 As η closer to 0: closer to uniform distribution (more conservative learning)

Simple Tuning: the Good News

• Worst-case optimal for $\eta = \sqrt{8 \ln(K)/T}$:

Regret
$$\leq \sqrt{T \ln(K)/2}$$

Proof idea:

- approximate our loss: $w_t \cdot \ell_t = \sum_k w_{t,k} \ell_{t,k}$

- by the **mix loss**:

$$m_t = \frac{-1}{\eta} \ln \sum_k w_{t,k} e^{-\eta \ell_{t,k}}$$

- and bound the approximation error:

$$\delta_t = w_t \cdot \ell_t - m_t$$

Simple Tuning: the Good News

our loss = mix loss + approx. error $w_t \cdot \ell_t = m_t + \delta_t$

• Cumulative mix loss is close to L*:

$$L^* \le \sum_{t=1}^{T} m_t \le L^* + \frac{\ln K}{\eta}$$

 $\delta_t \le \frac{\eta}{8}$

Hoeffding's bound:

• Together:

$$\sum_{t=1}^{T} w_t \cdot \ell_t - L^* \le \frac{\ln K}{\eta} + \frac{\eta T}{8} \xrightarrow{\eta = \sqrt{8 \ln K/T}} \sqrt{T \ln(K)/2}$$

Balances the two terms

Lost Advantages of FTL



 Simple tuning does much worse than FTL on i.i.d. losses

Simple Tuning: the Bad News

• The bad news:

- $-\eta = \sqrt{8\ln(K)/T}$ = conservative learning
- In practice, better when learning rate does not go to 0 with T! [DGGS, 2013]
- Lost advantages of FTL!

- We want to exploit luckiness:
 - robust against worst-case losses; but
 - if the data are `easy', we should learn faster!

Luckiness: Exploiting Easy Data

Improvement for small losses:

$$\mathsf{Regret} = O\left(\sqrt{L^* \ln(K)}\right)$$

variance of w_t

Second-order Bounds:

- [CBMS, 2007] and AdaHedge: $O\left(\sqrt{\sum_{t} v_t \ln(K)}\right)$
- Related bound by [HK, 2008]

Luckiness: Exploiting Easy Data

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$$O\left(\sqrt{\frac{L^*(T-L^*)}{T}\ln(K)}\right)$$

2nd-order Bounds: I.I.D. Data
variance of
$$w_t$$

Regret bound: $O\left(\sqrt{\sum_t v_t \ln(K)}\right)$

• For IID data, w_t concentrates fast on best expert:

$$\blacktriangleright \sum_{t} v_t \le C - \blacktriangleright \mathsf{Regret} \le C'$$

2nd-order Bounds: I.I.D. Data



Recover FTL benefits for i.i.d. data

CBMS: Proof Idea

our loss = mix loss + approx. error $w_t \cdot \ell_t = m_t + \delta_t$

- Cumulative mix loss is close to L^* : $L^* \leq \sum_{t=1}^T m_t \leq L^* + \frac{\ln K}{n}$
- Bernstein's bound:

$$\delta_t \leq \frac{1}{2}\eta v_t + \text{ lower order terms}$$

• Together:

balancing

 $\operatorname{\mathsf{Regret}} \leq \frac{\ln K}{\eta} + \frac{1}{2}\eta \sum_{t=1}^{T} v_t \longrightarrow O\left(\sqrt{\sum_t v_t \ln(K)}\right)$

AdaHedge: Proof Idea

our loss = mix loss + approx. error $w_t \cdot \ell_t = m_t + \delta_t$

• Cumulative mix loss is close to L*:

$$L^* \le \sum_{t=1}^T m_t \le L^* + \frac{\ln K}{\eta}$$

• No bound:

$$\delta_t = \delta_t$$

• Together: balancing $\eta = \frac{\ln(K)}{\sum_t \delta_t}$

 $\mathsf{Regret} \leq \frac{\ln K}{\eta} + \sum_{t} \delta_{t} \longrightarrow O\left(\sum_{t} \delta_{t}\right) = O\left(\sqrt{\sum_{t} v_{t} \ln K}\right)$

AdaHedge: Proof Idea

our loss = mix loss + approx. error $w_t \cdot \ell_t = m_t + \delta_t$

• Cumulative mix loss is close to L*:

$$L^* \le \sum_{t=1}^T m_t \le L^* + \frac{\ln K}{\eta}$$

 $\mathsf{Regret} \leq \frac{\ln K}{\eta} + \sum_{t} \delta_{t} \longrightarrow O\left(\sum_{t} \delta_{t}\right) = O\left(\sqrt{\sum_{t} v_{t} \ln K}\right)$

No bound:

$$\delta_t = \delta_t$$

NB Bernstein's bound is pretty sharp, so in practice CBMS ≈ AdaHedge up to constants.

• Together:

balancing

 $\eta = \frac{\ln(K)}{\sum_t \delta_t}$

Tuning η Online

- Balancing η in CBMS and AdaHedge depends on unknown quantities
- Solve this by changing $\eta = \eta_t$ with t
- Problem: $\sum_{t} m_t \leq L^* + \ln K/\eta$ breaks

Lemma [KV, 2005]: If $\eta_1 \ge \eta_2 \ge \eta_3 \ge \ldots$, then

$$\sum_{t=1}^{T} m_t \le L^* + \ln(K) / \boldsymbol{\eta_T}$$

2nd-order Bounds: the Bad News



 Do not recover FTL benefits for other `easy' data with a small number of leader changes

Luckiness: Exploiting Easy Data

Improvement for small losses:

$$\mathsf{Regret} = O\left(\sqrt{L^* \ln(K)}\right)$$

- Second-order Bounds:
 - [CBMS, 2007] and AdaHedge: $O\left(\sqrt{\sum_{t} v_t \ln(K)}\right)$
 - Related bound by [HK, 2008]

• FlipFlop:

- "Follow the leader if you can, Hedge if you must"
- Regret < best of AdaHedge and FTL</p>

FlipFlop

• FlipFlop bound:

$\textbf{Regret} \leq \begin{cases} 6 \cdot \text{FTL } \textbf{Regret} \\ 3 \cdot \text{AdaHedge } \textbf{Regret Bound} \end{cases}$

- Alternate Flip and Flop regimes
 - Flip: Tune $\eta_t = \infty$ like FTL
 - Flop: Tune η_t like AdaHedge
- (No restarts of the algorithm, like in `doubling trick'!)

FlipFlop: Proof Ideas

- Alternate Flip and Flop regimes
 - Flip: Tune $\eta_t = \infty$ like FTL
 - Flop: Tune η_t like AdaHedge

Analysing two regimes:

- 1. Relate mix loss for Flip to mix loss for Flop
- 2. Keep approximation errors balanced between regimes

1. Relating Mix Losses

 $\eta_1 \ge \eta_2 \ge \eta_3 > \dots$

• We violate condition of KV-lemma:

• But:

 $\sum_{t} m_t \le \sum_{t} m_t^{\text{flop}} + C \sum_{t \in \text{flop}} \delta_t$ $\leq L^* + \frac{\ln K}{\eta_T^{\text{flop}}} + C \sum_{t \in \text{flop}} \delta_t$ $=L^* + (C+1) \sum \delta_t$ t∈flop

2. Balance Approximation Errors

 Alternate regimes to keep approximation errors balanced:

$$\sum_{t \in \text{flip}} \delta_t \propto \sum_{t \in \text{flop}} \delta_t$$

$$\text{Regret} = \sum_t m_t - L^* + \sum_t \delta_t \le (C+2) \sum_{t \in \text{flop}} \delta_t + \sum_{t \in \text{flip}} \delta_t$$

$$\propto \begin{cases} \sum_{t \in \text{flip}} \sum_t \delta_t & \longrightarrow \text{FTL Bound} \\ \sum_{t \in \text{flop}} \sum_t \delta_t & \longrightarrow \text{AdaHedge Bound} \end{cases}$$

Small Nr Leader Changes Again



 FlipFlop exploits easy data, AdaHedge does not

FTL Worst-case Again



Summary

- Follow-the-Leader:
 - works well for `easy' data: i.i.d., few leader changes
 - but not robust to worst-case data

- Second-order bounds (e.g. CBMS, AdaHedge):
 - robust against worst case + can exploit i.i.d. data
 - but do not exploit few leader changes in general

FlipFlop: best of both worlds

Luckiness: What's Missing?

- FlipFlop:
 - "Follow the leader if you can, Hedge if you must"
 - Regret \leq best of AdaHedge and FTL
- But what if optimal η is in between AdaHedge and FTL?
- Can we compete with the best possible η chosen in hindsight?

References

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EXTRA SLIDES

No Need to Pre-process Losses

• Common assumption $\ell_{t,k} \in [0,1]$ requires translating and rescaling the losses

- CBMS:
 - Extension so this is not necessary.
 Important when range of losses is unknown!
- AdaHedge and FlipFlop:
 - Invariant under rescaling and translation of losses, so get this for free.

2nd-order Bounds: I.I.D. Data

- Regret bound: $O\left(\sqrt{\sum_t v_t \ln(K)}\right)$
- If w_t concentrates fast on best expert, then

$$\sum_{t} v_t \le C \longrightarrow \mathsf{Regret} \le C'$$

variance of w_t

• IID data:

1. Balancing $\eta_t = \sqrt{\frac{2 \ln(K)}{\sum_{s}^{t-1} v_s}}$ is large for all $t \leq T$ 2. w_t concentrates fast 3. Then 1. also holds for T + 1

FlipFlop on I.I.D. Data



Example: Spam Detection

	Subject	From	
	🛛 Gratis Turkije	Reizen Center	$y_1 = 1$
	uitnodiging hoorzitting reorganisatie FEW dinsdag.	20 se Ivo van Stokkum	$y_2 = 0$
	🖩 Re: Urgent Business Inquiry.	Ubc Ltd	$y_3 = 1$
	🖾 Reminder: first colloquium	Jeu, R.M.H. de	$y_4 = 0$
	@ Informatie over VUnet	College van Bestuur	$y_5 = 0$
	■ USD 500 Free Deposit at PartyPoker!	PartyPoker	$y_6 = 1$
		UK INTL. LOTTERY PROMOTION	$y_7 = 1$
	📾 bachelor/master diploma uitreiking 14 september	Sotiriou, M.	$y_8 = 0$
0	■ HAPPY NEW YEAR 2068	Anil Shilpakar	$y_9 = 1$
ê	a Thailand Package	Anil Shilpakar	$y_{10} = 1$

Example: Spam Detection

- Data: (x_t, y_t) with $y_t \in \{0, 1\}$
- Predictions: probability $p_t \in [0, 1]$ that $y_t = 1$
- Loss (probability of wrong label):

$$\ell(y_t, p_t) = \begin{cases} p_t & \text{if } y_t = 0\\ 1 - p_t & \text{if } y_t = 1 \end{cases}$$

- Experts: K spam detection algorithms
- If expert k predicts $p_{t,k}$, then $\ell_{t,k} = \ell(y_t, p_{t,k})$
- Regret: expected nr. mistakes over expected nr. of mistakes of best algorithm

FTL: the Bad News

- Consider two trivial spam detectors (experts): $p_{t,1} = 0$ $p_{t,2} = 1$
- If we deterministically choose an expert k*
 (like FTL) then we could be wrong all the time:

$$\ell_{t,k*} = 1 \quad \ell_{t,\neg k*} = 0$$

Regret:

- Let *n* denote the number of times expert 1 has loss 1. Then $L^* = \min\{n, T n\} \le T/2$
- Linear regret = $T L^* \ge T/2$

