

Switching Investments

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October 6, 2010



Centrum Wiskunde & Informatica

What We Do

All About A Line
Basic Investment
Strategies

Hedging

Price Switched
Strategies

More Price
Switching

What We Actually
Do

What We Do

All About A Line

What We Do

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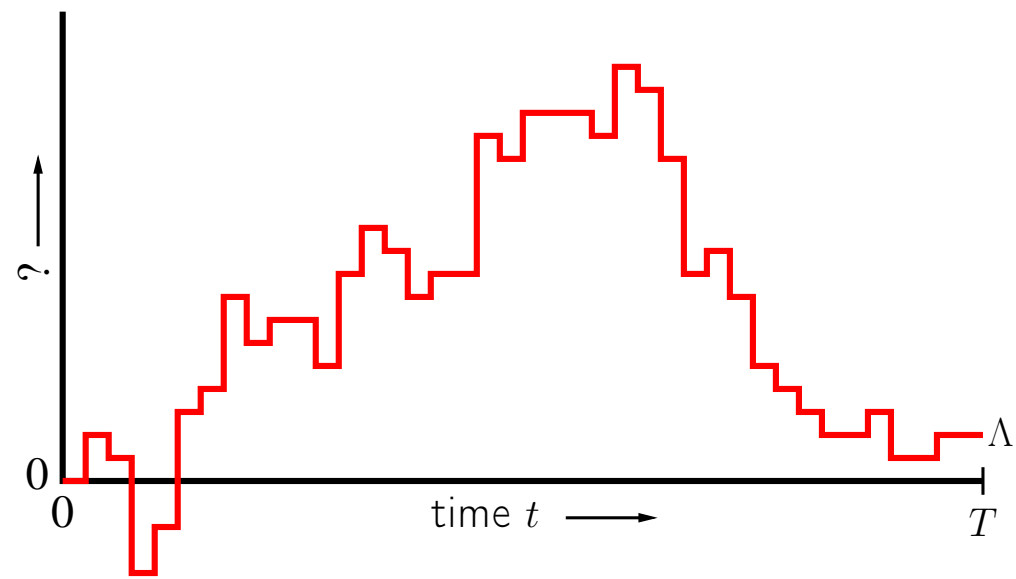
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Basic Investment Strategies

What We Do

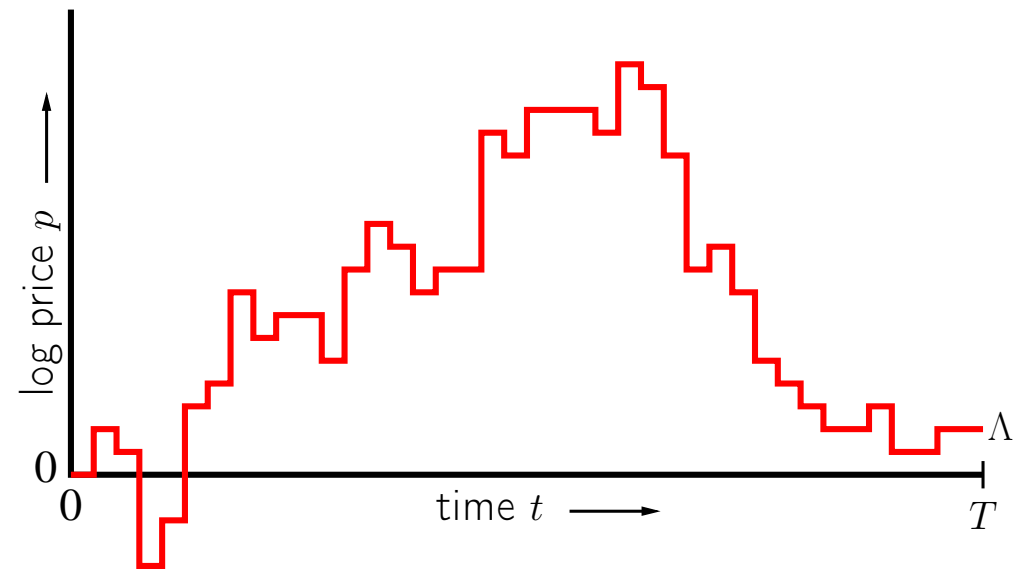
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A basic investment strategy σ_t is to sell at a predetermined time t .

Basic Investment Strategies

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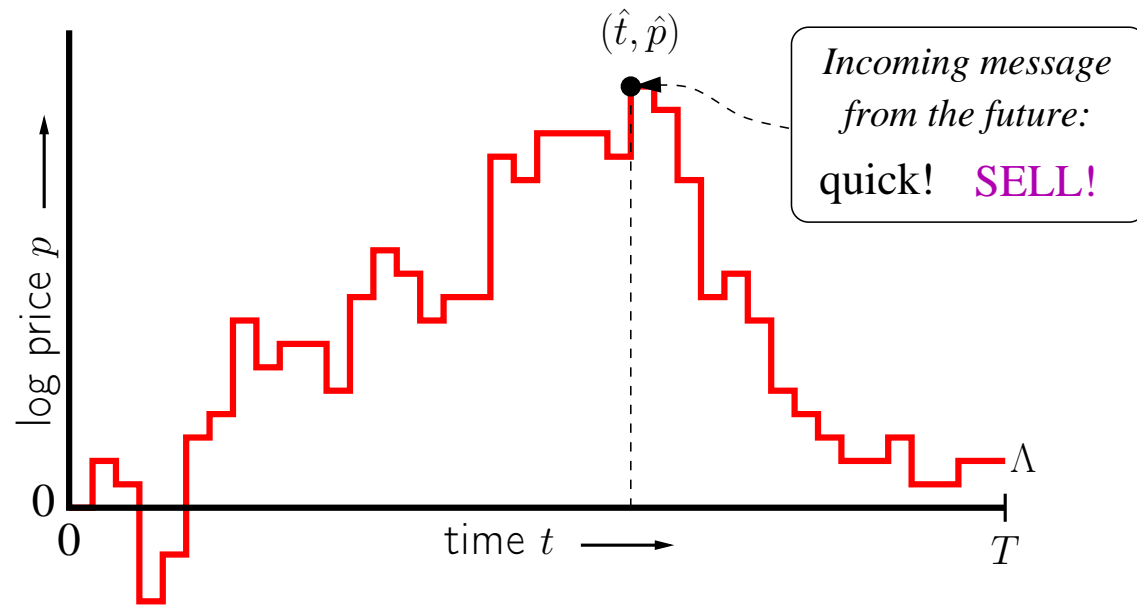
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A basic investment strategy σ_t is to sell at a predetermined time t .

Problem: *in hindsight* we know when the oil started leaking!

Hedging

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We distribute our initial capital \$1 over strategies $\sigma_0, \dots, \sigma_T$.

Let $\tau(t)$ denote the fraction of capital assigned to σ_t .

Let $\Lambda(0) = 0$. We obtain payoff:

$$\log \sum_{t=0}^T e^{\Lambda(t)} \tau(t) \geq \log \left(e^{\Lambda(\hat{t})} \tau(\hat{t}) \right) = \Lambda(\hat{t}) - (-\log \tau(\hat{t})).$$

Hedging

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Regret may be relatively large or small, depending on

- ✓ The granularity of measurement

Hedging

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Regret may be relatively large or small, depending on

- ✓ The granularity of measurement ← undesirable!

Price Switched Strategies

What We Do

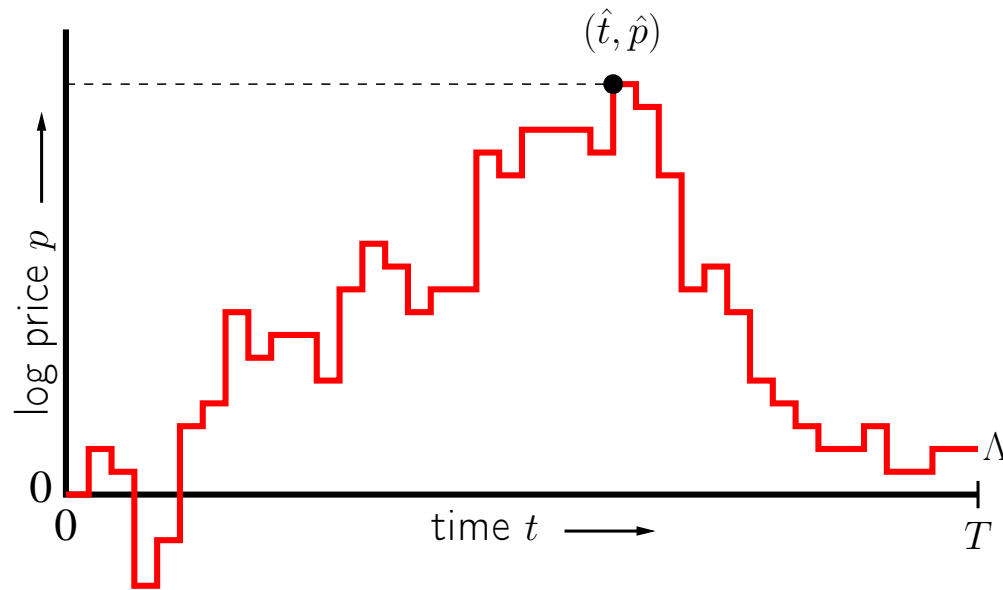
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Let us now define σ_p to sell when $\Lambda(t) \geq p$.

- ✓ Time-switched strategy σ_t : decision to sell depends on t
- ✓ Price-switched strategy σ_p : decision to sell depends on $\Lambda(t)$

Price Switched Strategies

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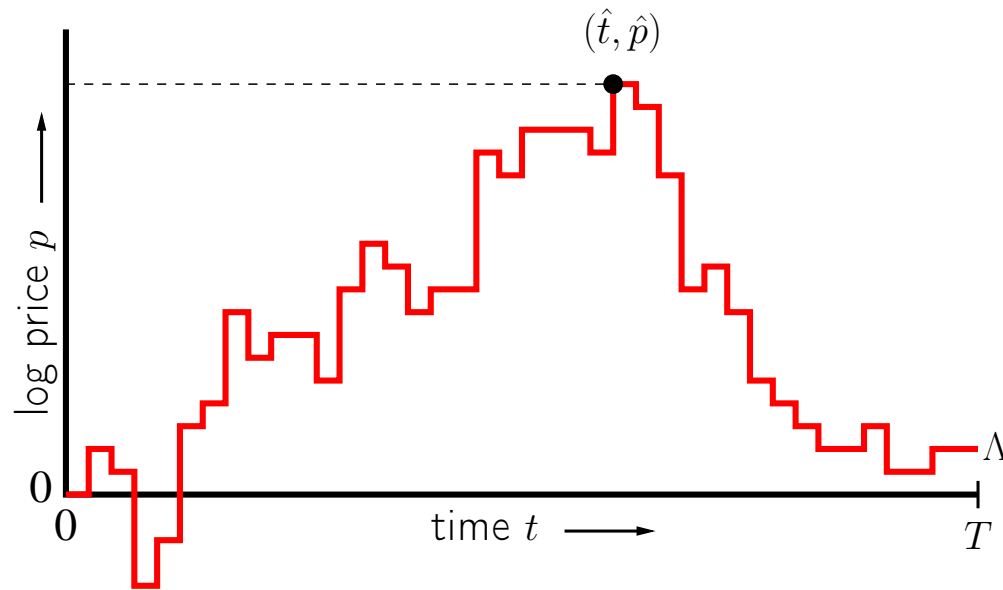
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Let us now define σ_p to sell when $\Lambda(t) \geq p$.

- ✓ Time-switched strategy σ_t : decision to sell depends on t
- ✓ Price-switched strategy σ_p : decision to sell depends on $\Lambda(t)$

We can no longer sell at every moment. But that's OK.

More Price Switching

What We Do

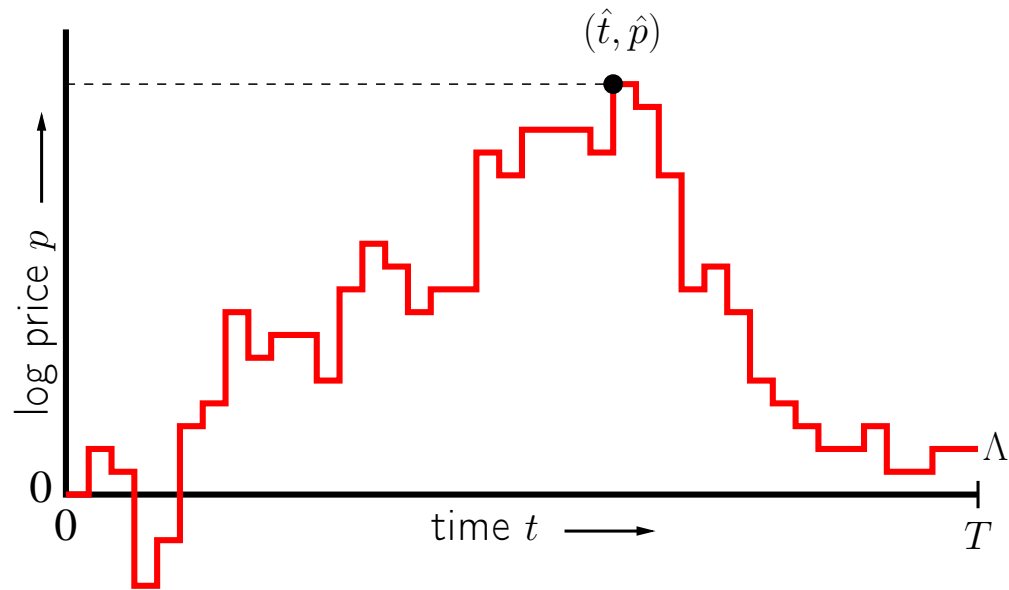
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We can hedge, now with π on price levels, to obtain at least

$$\log \sum_{p=0}^{\hat{p}} e^p \pi(p) \geq \log \left(e^{\hat{p}} \pi(\hat{p}) \right) = \underbrace{\hat{p}}_{\text{ideal}} - \underbrace{\left(-\log \pi(\hat{p}) \right)}_{\text{regret}}.$$

For sufficiently large \hat{p} , the regret is relatively small!

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Continuous Price
Multiple Switches
Continuous Time
Monotonicity
Regret Bound
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What We Actually Do

Continuous Price

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Actually, logprices are not integers and we do not pretend they are.

We can get very close to the previous bound:
if π is a decreasing density on the positive reals, then

$$\log \int_0^{\hat{p}} e^p \pi(p) dp \geq \log \left(\pi(\hat{p}) \int_0^{\hat{p}} e^p dp \right) = \underbrace{\log(e^{\hat{p}} - 1)}_{\approx \text{ideal } \hat{p}} - \underbrace{(-\log \pi(\hat{p}))}_{\text{regret}}.$$

We cannot sell at \hat{p} exactly anymore \rightarrow small additional overhead

Multiple Switches

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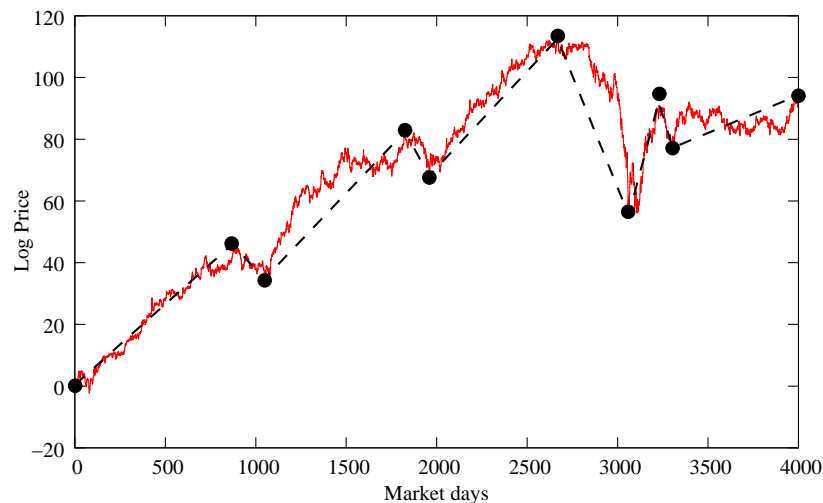
Example

Algorithm

Actually, we are interested in exploiting multiple switches.

Let $\delta = (\delta_1, \delta_2, \dots)$. A strategy σ_δ :

- ✓ initially invests all capital
- ✓ sells all stock when the logprice goes up δ_1 or more, then
- ✓ invests all capital again as it goes down δ_2 or more,
- ✓ etcetera.



To hedge, take the infinite product distribution of π .

Continuous Time (Theorem 1)

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Intuition: Discontinuities in Λ are helpful.

Continuous Time (Theorem 1)

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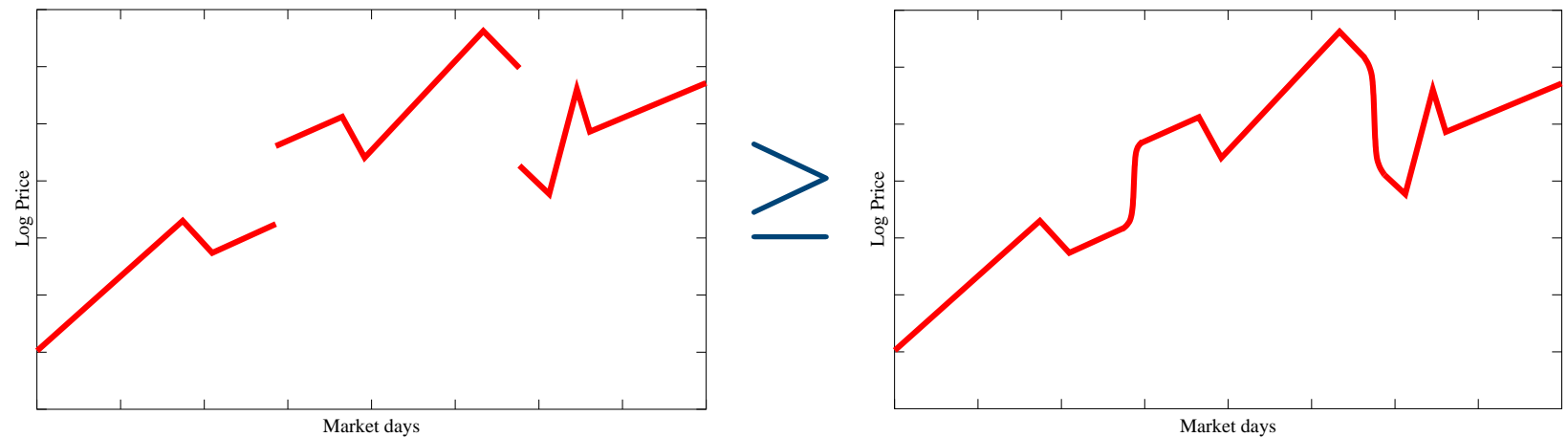
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Intuition: Discontinuities in Λ are helpful.

Let the logprice function be $\Lambda : [0, T] \rightarrow \mathbb{R}$.

(A discrete time scenario can be modelled by a step function.)



Continuous Time (Theorem 1)

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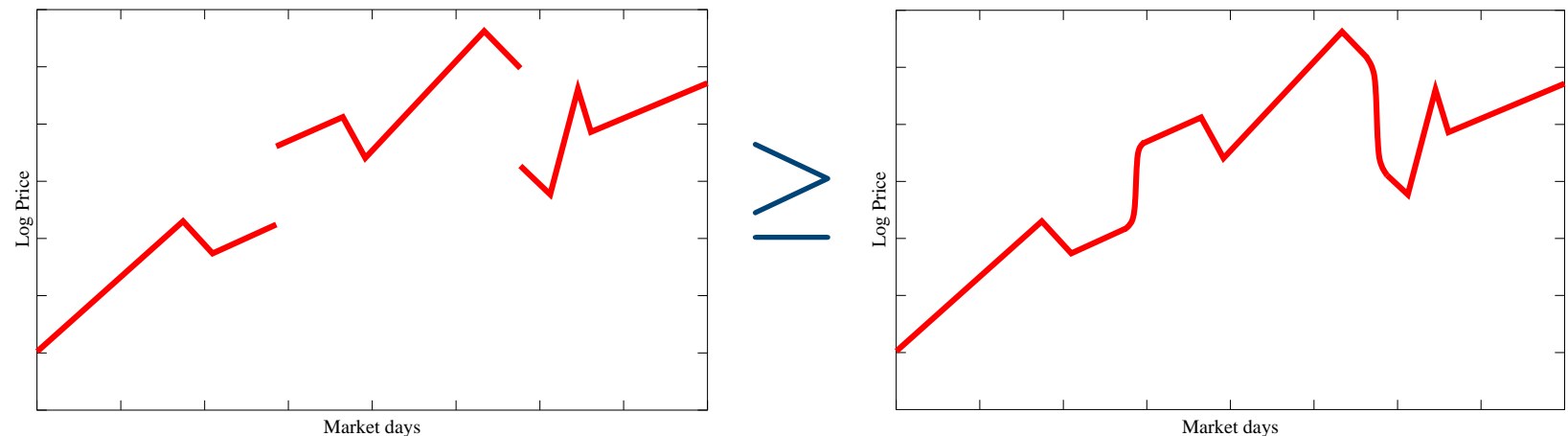
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Let the logprice function be $\Lambda : [0, T] \rightarrow \mathbb{R}$.

(A discrete time scenario can be modelled by a step function.)



✓ We can simplify the analysis by assuming continuity.

Monotonicity (Theorem 2)

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Intuition: The more fluctuations in Λ , the better.

Monotonicity (Theorem 2)

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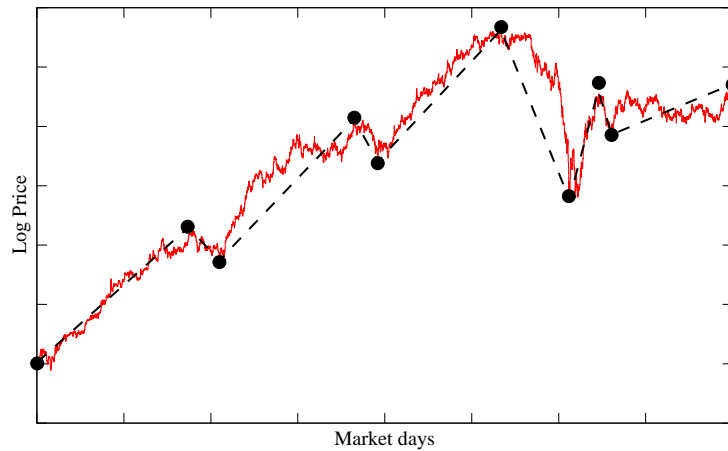
Monotonicity

Regret Bound

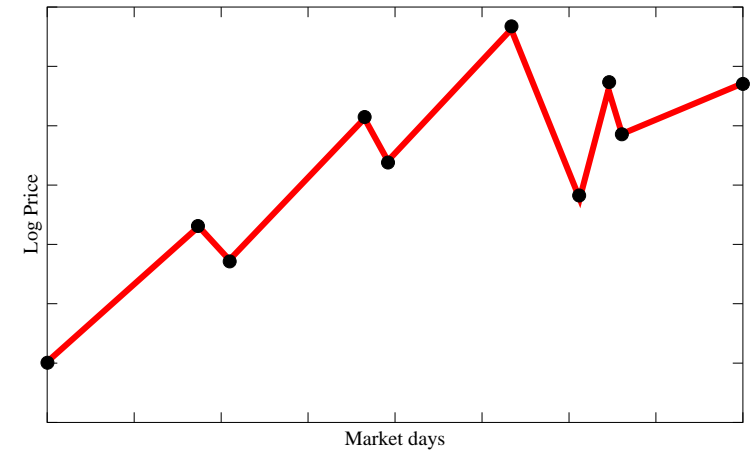
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\geq



Monotonicity (Theorem 2)

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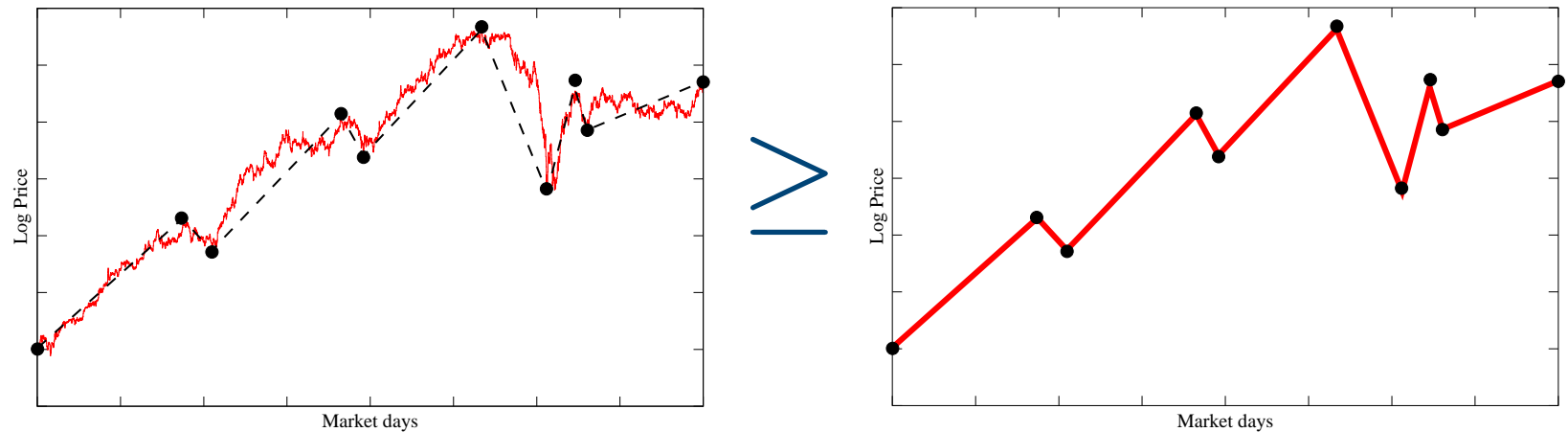
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Intuition: The more fluctuations in Λ , the better.



In summary, the regret compared to a specific σ_δ is maximised if

- ✓ Λ is continuous (Thm 1)
- ✓ Λ is monotonic in-between switches (Thm 2)

The worst case for regret coincides with the ideal case for analysis!

Regret Bound

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Theorem 3 Fix Λ . For any basic strategy σ_δ that performs its m^{th} switch on Λ at time T , the payoff of our strategy is at least

$$\underbrace{\sum_{1 \leq \text{odd } i \leq m} \delta_i}_{\text{ideal}} - \underbrace{\sum_{i=1}^m -\log \pi(\delta_i)}_{\text{regret}} - m \cdot \text{small}.$$

Regret Bound

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Thus,

- ✓ Small fluctuations are hard to exploit
- ✓ The bound is best applied to parsimonious strategies (with small m)

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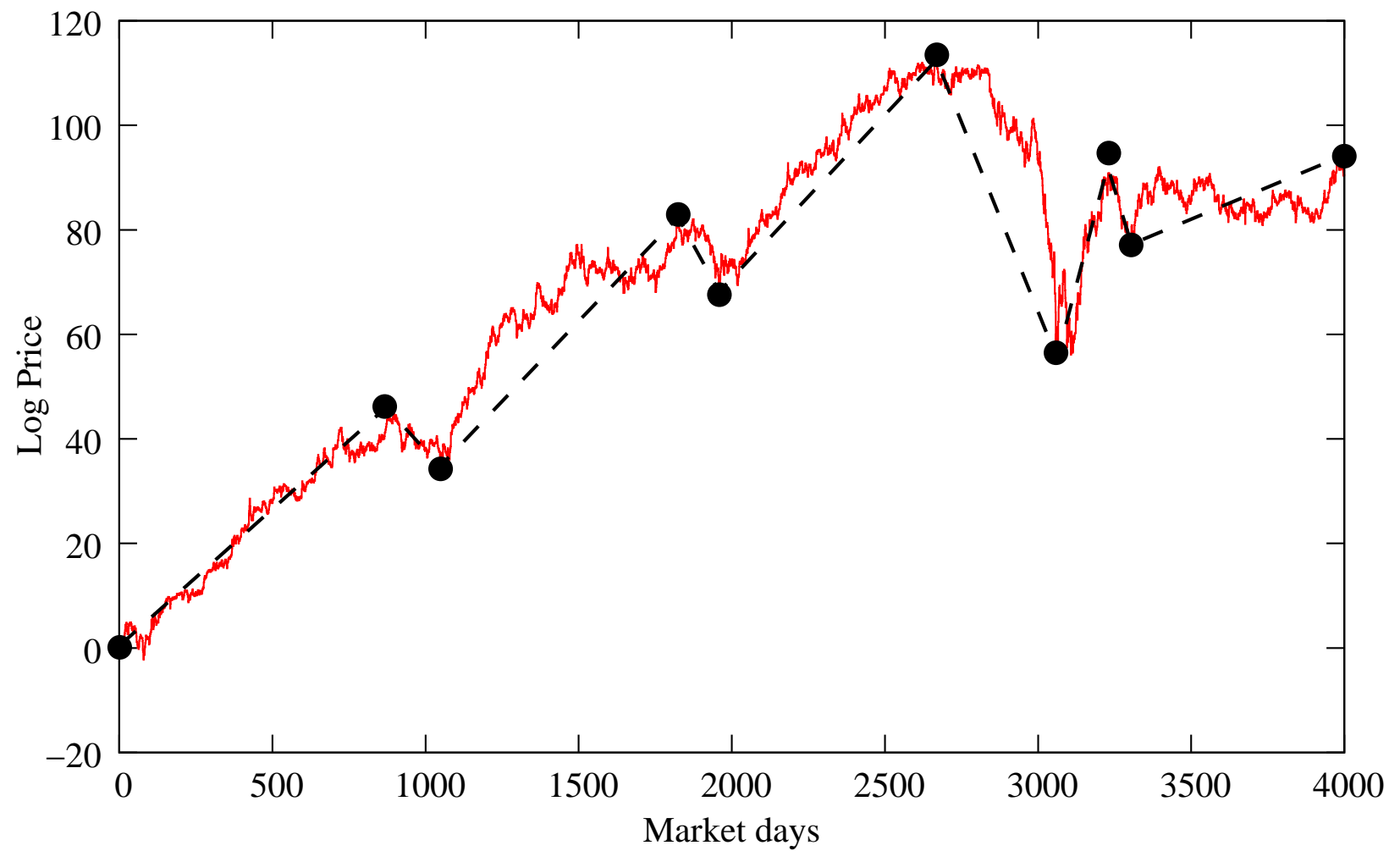
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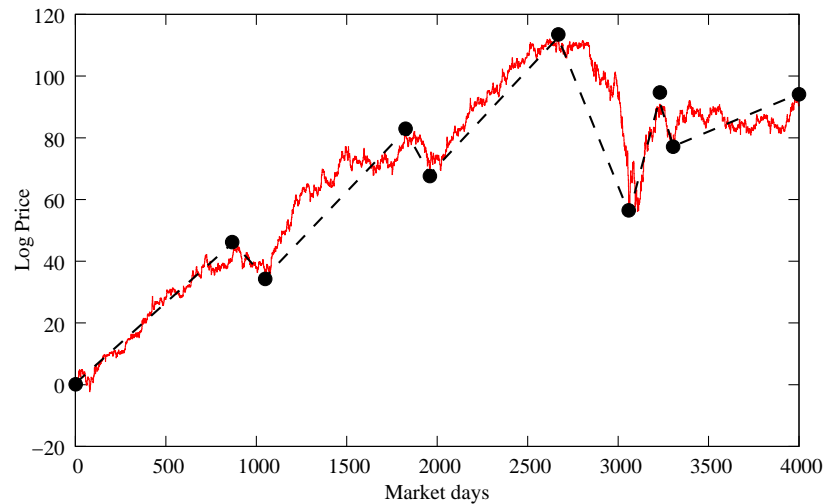
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Example
Algorithm



Strategy	Payoff
Invest everything	90
Ideal	1021
Model	178
Bound	105
Actual performance	175

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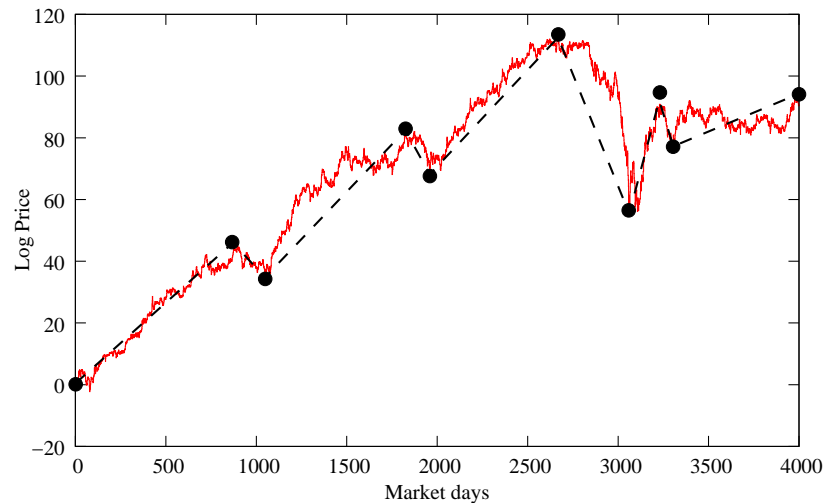
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- ✓ Performance on real stock: probably not brilliant
- ✓ Strategy still useful as a safeguard against excessive loss

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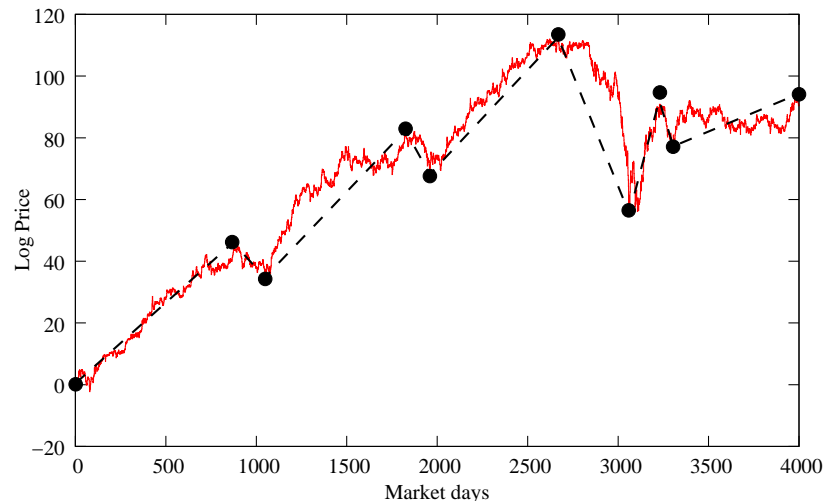
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- ✓ Performance on real stock: probably not brilliant
- ✓ Strategy still useful as a safeguard against excessive loss
- ✓ In other applications Λ is usually less adversarial
- ✓ Performance is competitive with Fixed Share and typically better than Variable Share for log loss.

Algorithm

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Algorithm

A simple algorithm is described in the paper:

- ✓ Statisticians: “It’s just Bayes”
- ✓ Learning Theorists: “It’s just the Aggregating Algorithm”

Algorithm

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- ✓ If π is memoryless (exponential) running time can be reduced to $O(n)$.

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- ✓ Runs in $O(n^2)$ time and $O(n)$ memory.
- ✓ If π is memoryless (exponential) running time can be reduced to $O(n)$.
- ✓ It buys when you’re losing, and sells when you’re winning?!

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Thanks