

Adaptive Hedge

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SUMMARY

AdaHedge is a new online learning algorithm that adapts to the difficulty of the data

Difficulty	Regret
Worst-case data	$O(\sqrt{L_T^* \ln(K)})$
Easy data	constant: $O(K)$

Key Ideas

- Bounds on the *mixability gap* (see top-right panel) play a crucial role in previous analyses of the Hedge algorithm.
- We only bound the mixability gap in the analysis, but not in the algorithm!
- On easy data, the probabilities output by Hedge converge on a single action. In this case we improve the standard bounds.
- Example: if one action is always better than all others.

ONLINE LEARNING SETTING

Decision Theoretic Online Learning

In rounds $t = 1, \dots, T$:

1. Assign probabilities $w_t = (w_t^1, \dots, w_t^K)$ to K actions
2. Actions get losses $\ell_t \in [0, 1]^K$
3. Our loss: $w_t \cdot \ell_t$

Aim to minimize the *regret*

$$R(T) = \sum_{t=1}^T w_t \cdot \ell_t - L_T^*,$$

where $L_T^* = \min_k \sum_{t=1}^T \ell_t^k$ is the loss of the best action in hindsight.

HEDGE

- Hedge predicts with exponential weights:

$$w_t^k \propto \exp\left(-\eta \sum_{s=1}^{t-1} \ell_s^k\right).$$

- Its performance depends strongly on the *learning rate* $\eta > 0$.

MIXABILITY GAP

The *mixability gap* is

$$\delta_t(\eta) = w_t \cdot \ell_t - \left(-\frac{1}{\eta} \ln(w_t \cdot e^{-\eta \ell_t})\right).$$

- In Prediction with Expert Advice terms: $\delta_t(\eta)$ measures the difference with a mixable loss function.
- In Bayesian terms: $\delta_t(\eta)$ measures the difference between randomizing according to the posterior and mixing according to the posterior.

ADAHEDGE

- Tune η optimally for a budget $b(\eta)$ on the cumulative mixability gap $\Delta_T(\eta) = \sum_{t=1}^T \delta_t(\eta)$
- Increase the budget using the doubling trick.

Algorithm

1. Start with $\eta = 1$
2. Run a new instance of Hedge with learning rate η until $\Delta_T(\eta)$ exceeds budget

$$b(\eta) = \left(\frac{1}{\eta} + \frac{1}{e-1}\right) \ln(K).$$
3. Set $\eta \leftarrow \eta/2$ and goto 2.

THEORETICAL RESULTS

AdaHedge is worst-case optimal...

Theorem 1 The regret of AdaHedge is bounded by

$$R(T) \leq 5.1 \sqrt{L_T^* \ln(K)} + O(\ln(L_T^* + 2) \ln(K)).$$

...and has strong theoretical guarantees on 'easy' data

Theorem 2 Suppose the loss vectors ℓ_t are independent random variables and there exists a k^* such that

$$\min_{k \neq k^*} \mathbb{E}[\ell_t^k - \ell_t^{k^*}] > 0 \quad \text{for all } t \in \mathbb{Z}^+.$$

Then with probability at least $1 - \delta$ the regret of AdaHedge is bounded by a constant:

$$R(T) = O(K + \log(1/\delta)).$$

PROOF TECHNIQUES

Everyone bounds the mixability gap δ_t .

Standard Analysis

- Optimize η after bounding $\delta_t(\eta) \leq \eta/8$.

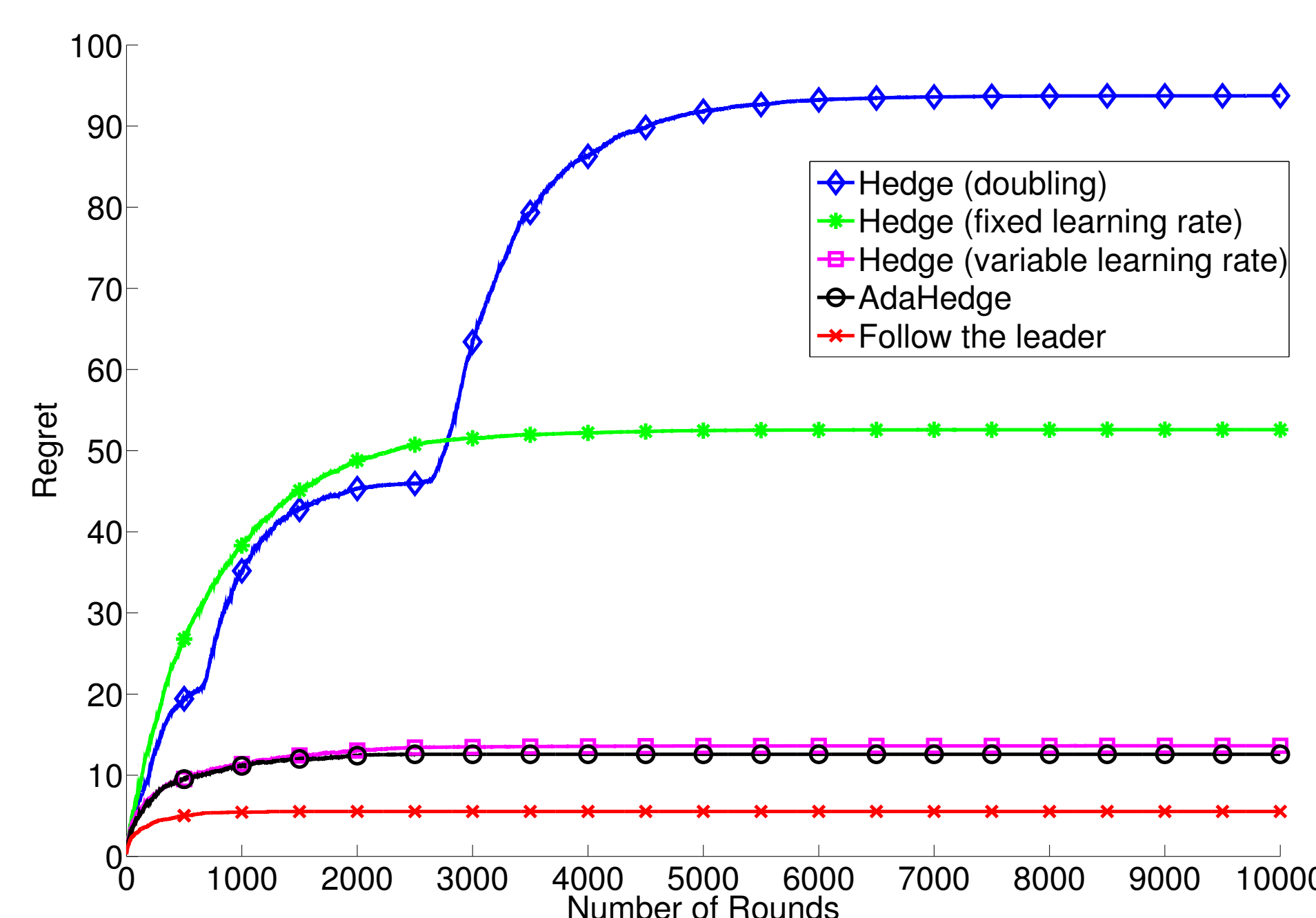
Our Approach

- Optimize η before bounding!
- If the posterior probabilities w_t converge on a single action, the mixability gap goes to 0!

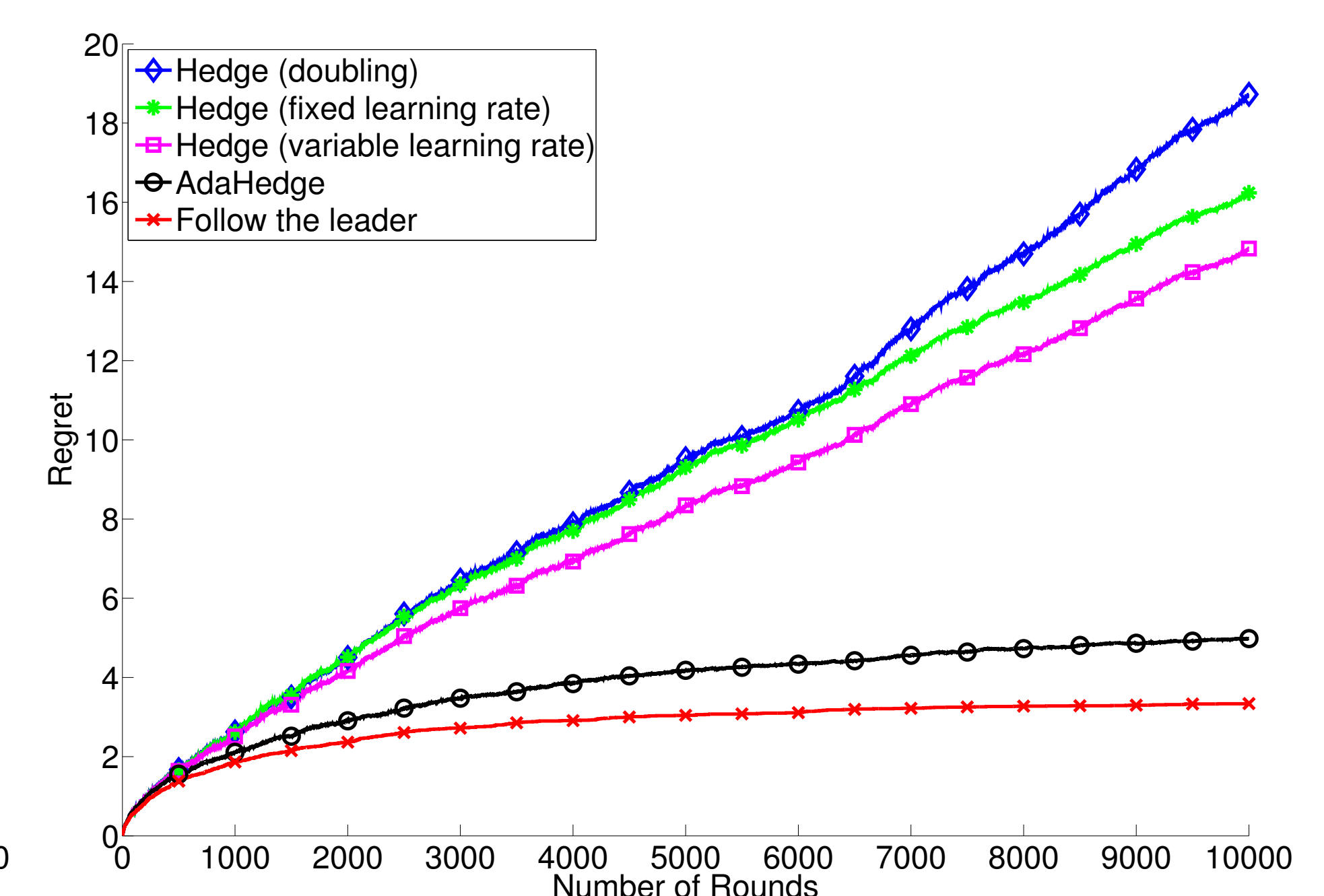
$$\delta_t(\eta) \leq (e-2)\eta(1 - \max_k w_t^k) \quad (0 < \eta \leq 1)$$

EXPERIMENTS

Simulation Study on 'Easy' Data



I.I.D. losses



Correlated losses

AdaHedge has excellent practical performance

N.B. Follow-the-leader does very well here, but gets *linear* regret $\geq T/2 - 1$ in the worst case!

CURRENT WORK

Avoid the Doubling Trick

- Better performance in practice
- Still very clean analysis
- Improved the worst-case bound to

$$R(T) \leq 2\sqrt{\frac{L_T^*(T - L_T^*)}{T} \ln(K)} + \frac{8}{3} \ln(K) + 2.$$

Weaker Conditions for Easy Data

- Guarantee regret bounded by the best *regret* of AdaHedge and Follow-the-Leader, up to a small constant factor.

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